

# Accelerating Full Configuration Interaction Quantum Monte Carlo

William Vigor <sup>1</sup>, James Spencer <sup>2,3</sup>, Michael Bearpark <sup>1</sup>,  
Alex Thom <sup>1,4</sup>

<sup>1</sup>Department of Chemistry, Imperial College London

<sup>2</sup>Department of Physics, Imperial College London

<sup>3</sup>Department of Materials, Imperial College London

<sup>4</sup>Department of Chemistry, University of Cambridge

# Part I

## Introduction: The Full Configuration Interaction Quantum Monte Carlo Algorithm

# Solving the Schrödinger Equation in a Basis of Slater Determinants

- To estimate the ground state energy of a Molecule with  $N$  electrons solve:

$$\hat{H}|\Psi\rangle = E|\Psi\rangle \quad (1)$$

where  $|\Psi\rangle$  is in some finite basis (at fixed geometry).

- If we let  $|\Psi\rangle$  be a single determinant of  $M$  Molecular Spin Orbitals (MSO), Hartree-Fock lets us compute the variations minimal single determinant wavefunction  $|D_0\rangle$ .
- Full Configuration Interaction (FCI) adds into  $|\Psi\rangle$  all possible determinants  $|D_I\rangle$  with  $N$  MSO occupied and  $M - N$  unoccupied .

# Full Configuration Interaction

- Want to solve the eigenvalue problem  $H|\Psi\rangle = e|\Psi\rangle$  where  $|\Psi\rangle = \sum_I C_I |D_I\rangle$
- Cast in to a matrix problem diagonalise matrix of  $\langle D_I | \hat{H} | D_J \rangle$
- Unfortunately this turns out to be impossibly large:  $\binom{M}{N}$  by  $\binom{M}{N}$ .
- **Can use Monte Carlo to sample this space.**

G. H. Booth, A. J. W. Thom, A. Alavi, J. Chem. Phys. **131**, 054106 (2009)

# Full Configuration Interaction Quantum Monte Carlo (FCIQMC) vs FCI

- Little communication overhead, can use massively parallel computers.
- Only need to store a stochastic representation of the eigenvector.
- Stochastic, quantifiable error (system and dynamics dependent) can be reduced to FCI accuracy by running for longer.
- **How quickly does this error converge?**
- **Can we find most efficient FCIQMC algorithm? i.e. how can we make the stochastic error converge quickest as a function of computer time?**

# Full Configuration Interaction Quantum Monte Carlo: Projector

- Denote the exact solutions (in a finite basis):

$$|\Psi_0\rangle, |\Psi_1\rangle, \dots, |\Psi_N\rangle \quad (2)$$

with energies:

$$E_0, E_1, \dots, E_N \quad (3)$$

We can expand our wavefunction in this basis:

$$|\psi\rangle = \sum_i^N a_i(\tau) |\Psi_i\rangle \quad (4)$$

If we apply the diffusion equation.

$$\frac{\partial |\psi\rangle}{\partial \tau} = -(\hat{H} - E_r \hat{1}) |\psi\rangle \quad (5)$$

and take  $\langle \Psi_j |$ :

$$\frac{\partial a_j(\tau)}{\partial \tau} = -(E_j - E_r) a_j(\tau) \quad (6)$$

Then:

$$a_j(\tau) = e^{-(E_j - E_r)\tau} a_j(\tau = 0) \quad (7)$$

Excited states die as  $\tau \rightarrow \infty$  and if  $E_r < E_0$  the  $|\Psi_0\rangle$  contribution grows.

# Full Configuration Interaction Quantum Monte Carlo: Discretised Determinant Space

First discretise our determinant space:

- Positive coefficients are represented by a number of positive psips ( $\Psi$  particles).
- Negative coefficients are represented by a number of negative psips.

$$|\Psi(\tau)\rangle = +0.25 |D_0\rangle + 0.5 |D_1\rangle - 0.25 |D_2\rangle$$

J . B. Anderson J. Chem. Phys. **63**, 1499 (1975)

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$$|\Psi(\tau)\rangle = \text{green circle} |D_0\rangle + \text{green circle} |D_1\rangle + \text{red circle} |D_2\rangle$$

J . B. Anderson J. Chem. Phys. **63**, 1499 (1975)



# Full Configuration Interaction Quantum Monte Carlo

- Can move  $\tau$  forwards by  $\delta\tau$  by applying:

$$e^{-(\hat{H}-E_r\hat{1})\delta\tau} \approx (\hat{1} - (\hat{H} - E_r\hat{1})\delta\tau) \quad (8)$$

- FCIQMC stochastically applies Eq. 8 in 3 steps, on each iteration.
  - 1 Spawn
  - 2 Diagonal Death
  - 3 Annihilation
- These define the psip population dynamics of FCIQMC.
- Thus the ground state wavefunction comes out if we start with a  $\psi$  such that  $\langle\psi|\hat{H}|\Psi\rangle \neq 0$  and run for many steps.

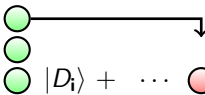
# FCIQMC Population Dynamics 1. Spawning

- Each psip attempts to spawn a child psip on a randomly selected determinant.

$$|\Psi(\tau)\rangle = \begin{array}{c} \circ \\ \circ \end{array} |D_0\rangle + \cdots \begin{array}{c} \circ \\ \circ \\ \circ \end{array} |D_i\rangle + \cdots \circ |D_j\rangle + \cdots$$

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- with probability  $(\langle D_i | H | D_j \rangle) \delta\tau$
- if  $(\langle D_i | H | D_j \rangle) \delta\tau < 0$  psip has same sign as parent and vice-versa.

# FCIQMC Population Dynamics 2. Diagonal Death

- Each parent psip attempts to die or is cloned.

$$|\Psi(\tau)\rangle = \begin{array}{c} \circ \\ \circ \end{array} |D_0\rangle + \dots \begin{array}{c} \circ \\ \circ \\ \circ \end{array} |D_i\rangle + \dots \circ |D_j\rangle + \dots$$

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- Death occurs with probability  $(\langle D_i|H|D_i\rangle - E_r)\delta\tau$
- Cloning occurs (population becomes more negative or positive) if  $(\langle D_i|H|D_i\rangle - E_r) > 0$
- $E_r$  is initially set to the energy of the Hartree-Fock until the psip population reaches the desired level.
- After which  $E_r$  is periodically updated to keep the population at the desired level.

# FCIQMC Population Dynamics 3. Annihilation

- Positive and negative psips residing on the same determinant annihilate:

$$|\Psi(\tau)\rangle = \begin{array}{c} \circ \\ \circ \end{array} |D_0\rangle + \cdots \begin{array}{c} \circ \\ \circ \\ \circ \end{array} |D_i\rangle + \cdots \begin{array}{c} \circ \\ \circ \end{array} |D_j\rangle + \cdots$$

J. S. Spencer, N. S. Blunt, and W. M. C. Foulkes J. Chem. Phys. **136**, 054110 (2012)

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- Essential so that the step-wise average of psip vector is the eigenvector with smallest eigenvalue of the FCI matrix.
- If we use enough psips.

J. S. Spencer, N. S. Blunt, and W. M. C. Foulkes J. Chem. Phys. **136**, 054110 (2012)

# Accumulating Expectation Values

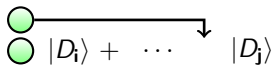
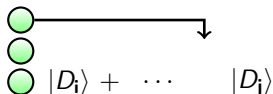
- Projected energy:

$$E = \frac{\langle D_0 | \hat{H} e^{\hat{H}\tau} | D_0 \rangle}{\langle D_0 | e^{\hat{H}\tau} | D_0 \rangle} = \frac{\langle D_0 | \hat{H} | \Psi_0 \rangle}{\langle D_0 | \Psi_0 \rangle} \quad (9)$$

- $\langle D_0 | \hat{H} e^{\hat{H}\tau} | D_0 \rangle$  equals a factor of:
  - $\langle D_0 | H | D_i \rangle$
  - $-\langle D_0 | H | D_i \rangle$  sum over every psip, sum over every determinant.
- $\langle D_0 | \Psi_0 \rangle$  equals the number of psips on the Hartree–Fock.
- Average over every step after the simulation has equilibrated.
- Need to estimate errors carefully due to serial correlation. As  $\delta\tau$  is small the stochastic representation of the eigenvector changes only a small amount between iterations, correlating estimates close in iteration space.

# The Initiator Approximation

- Allow only psips on determinants with a population above a threshold (initiator determinants) to spawn onto unoccupied determinants.

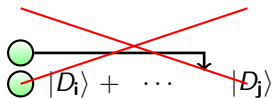
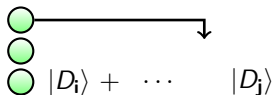


- Error bars converge faster and we need fewer psips.

D. Cleland, G. H. Booth, A. Alavi J. Chem. Phys. **132**, 041103 (2010)

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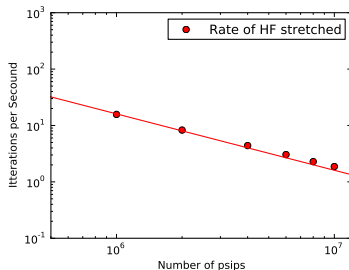
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## Part II

How can we use current algorithms most effectively

# Computer time as a function of the number of psips

- In FCIQMC more psips equivalent to more steps in terms of computer time.
- $\mathcal{O}(l \log l)$  sorting of newly spawned psips before of annihilation is negligible ( $l$  length of list to sort).
- Fits to  $\frac{c}{N_p}$  shown ( $N_p$  number of psips).



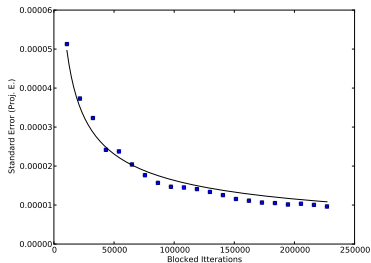
- What about the error bar?
- Is it better to run for more steps or use more psips (which is a better use of resources)?

# Quantifying Efficiency Empirically 1

- For any Monte Carlo algorithm the stochastic error (in some expectation value of the simulation) as a function of the number of steps  $N$  :

$$\sigma = \frac{a}{\sqrt{N}} \quad (10)$$

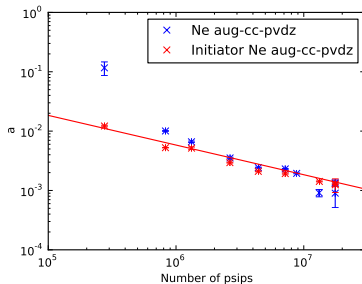
- So can use  $a$  to quantify the efficiency of FCIQMC, provided no systematic error is introduced.
- Small  $a$  is good, errors converge fast.



# Empirical Results: Scaling with the number of psips 1

$$\sigma = \frac{a}{\sqrt{N_s}} \quad (11)$$

- How does  $a$  depends on the number of psips  $N_p$  in the simulation ( $N_s$  number of iterations).
- One should fill the memory with psips if  $a \rightarrow 0$  faster than  $\mathcal{O}(\frac{1}{\sqrt{N_p}})$ .



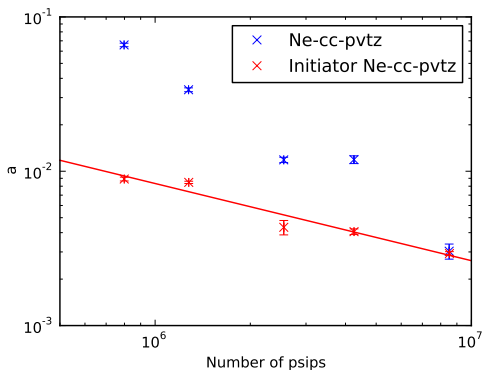
- For FCIQMC  $a$  decays faster than  $\mathcal{O}(\frac{1}{\sqrt{N_p}})$ , for i-FCIQMC seems to decay as  $b\sqrt{N_p}$  meaning more psips equivalent to more iterations.



# Empirical Results: Scaling with the number of psips 2

$$\sigma = \frac{a}{\sqrt{N_s}} \quad (12)$$

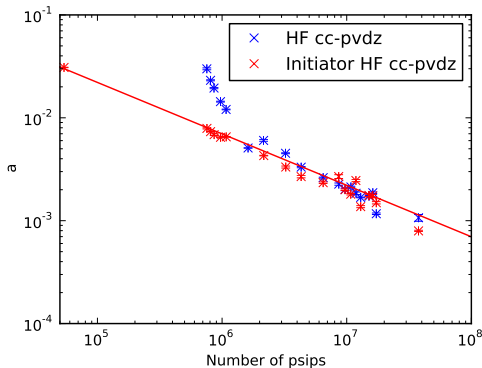
- Other systems show similar results.



# Empirical Results: Scaling with the number of psips 3

$$\sigma = \frac{a}{\sqrt{N_s}} \quad (13)$$

- Other systems show similar results.



# Conclusions

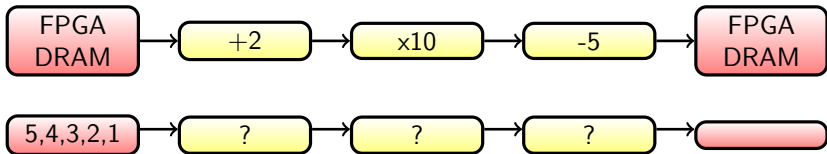
- For calculations with the initiator approximation more psips is equivalent to more steps in terms of  $\sigma(E)$ .
- For other systems the error decreases faster as the number of psips increase getting closer to the initiator limit.
- At the limit of a large number of psips, trivial parallelism mode (running multiple independent simulations).
- We now have a way of comparing different Monte Carlo algorithms.

# Part III

## Novel Hardware

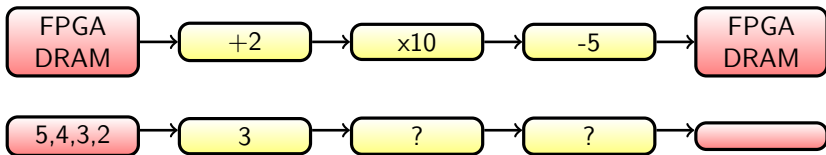
# Field Programmable Gate Arrays (FPGA's)

- A programmable grid of logic components.
- Components can be grouped and connected so that each group performs an operation and passes it too the next group each clock cycle.
- Sending data back up the chip can be tricky as we have to meet timings.
- Clock rate of  $\sim 100$  MHz c.f. CPU  $\sim 1000$  MHz .



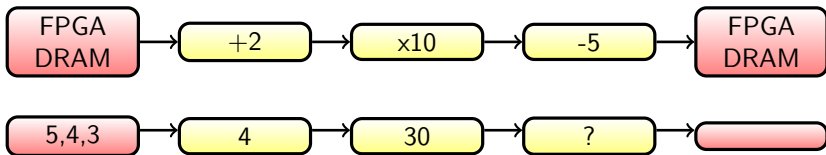
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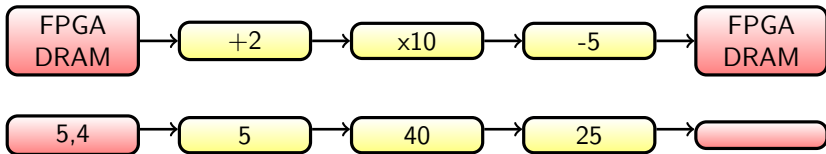
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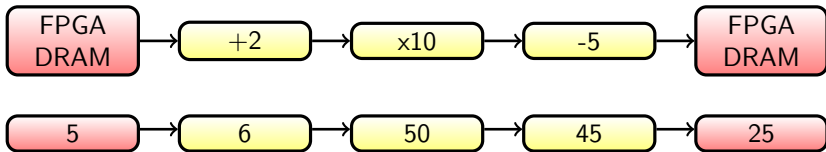
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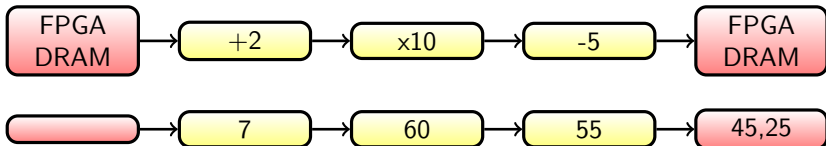
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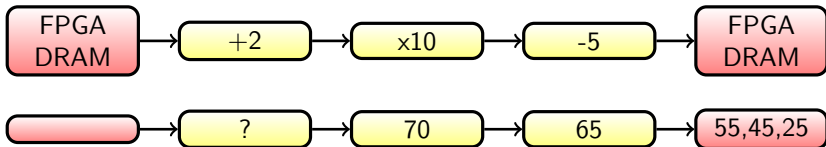
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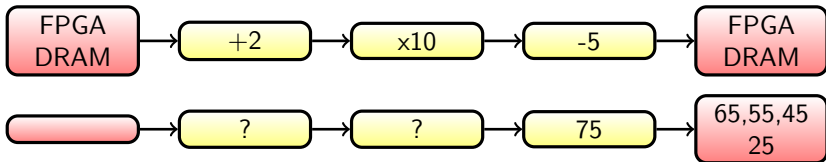
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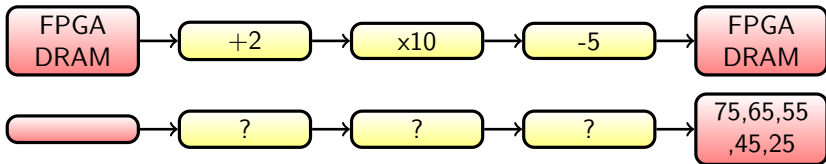
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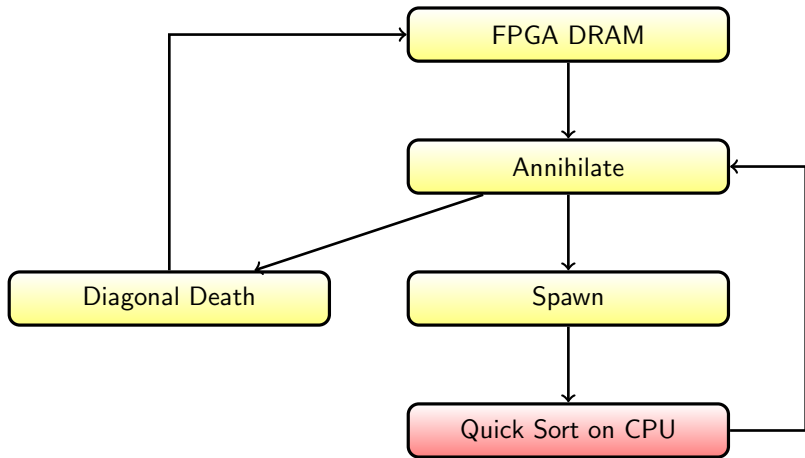
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- The psip vector makes a good candidate for streaming through the FPGA.
- Diagonal death  $\mathcal{O}(L)$  and spawning  $\mathcal{O}(N_p)$ . Where  $L$  is the length of psip vector.
- Annihilation requires sorting the newly spawned psips  $\mathcal{O}(l \log l)$
- Efficient sorting tricky.
- Send to CPU for sort.
- Currently investigating the Hubbard model (can store the integrals on chip as they have a simple structure).

# FCIQMC on FPGA's



# Conclusions and Directions

- Move to multiple FPGAs and CPUs.
- Investigate the Hubbard Model and quantify speed up.
- Real chemical systems require access of the integrals to generate the matrix elements, this will be tricky.
- It may be more efficient for all psips on a determinant to spawn at once Diagonal Death and spawning both  $\mathcal{O}(L)$ .
- Is this version more efficient for an FPGA ? (Work in progress)



# Acknowledgements

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- EPSRC for a studentship.
- Stephen Girdlestone, Craig Davies, and Robin Bruce at Maxeler Technologies
- All calculations ran using HANDE (developed at Imperial College) and the Imperial College High Performance Computing Service.



# Serial Correlation

- As  $\tau$  is small the vector of psips only changes a small amount between iterations.
- Thus expectation values close in iteration space are correlated.
- One has to remove this serial correlation by blocking: Average into blocks and compute standard error of blocks.
- We used the Iterative algorithms for optimal block size in:  
R. M. Lee, G. J. Conduit, N. Nemec, P. López Ríos, N. D. Drummond, Phys. Rev. E **83**, 066706 (2011)

